## Symbolic Robustness Analysis

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Presented by: Sayali Salvi, Saarland University Seminar on "Robustness of Hardware and Software Systems"

## Outline

- Control Systems
- Robustness
- Symbolic Robustness Analysis
- Algorithm and Implementation
- Limitations

## **Control Systems**



### Robustness

Robust



Question: Is the implementation still Robust?

• We say system is robust, when

small perturbations in the system inputs cause only small changes in its outputs.

### Robustness



System is  $(\delta, \varepsilon)$ -robust in input  $x_1$ 

## **Transmission Calculation Example**

int calc trans slow torques (int angle, int speed) data\_table: { int pressure1, pressure2; int gear, val1, val3, val4; val1=41  $val1 = lookup1 (\&(data_table[0][0]), angle);$ out1 table: if  $(3 * speed \le val1)$ gear = 3;gear=4 else gear = 4; $val3 = lookup2 (\&(out1_table[0][0]), gear);$ out2 table: pressure1 = val3 \* 1000;  $val4 = lookup2 (\&(out2_table[0][0]), gear);$ val4=1pressure2 = val4 \* 1000;angle = 30, speed = 13angle = 30, speed = 14pressure2 = 0pressure2 = 1000

angle	val1
30	41
40	63

gear	val3
1	0
2	0
3	1
4	1

gear	val4
1	0
2	0
3	0
4	1

## **Robustness Analysis**

- Why is it difficult?
  - huge input space to be tested exhaustively
  - many different code execution paths
  - many control computations based on table lookups
- Why is it required?
  - Random testing ineffective

 This paper studies robustness analysis for control software: using 'Symbolic Execution'

## Symbolic Execution

• Test generation technique

- Executes program on symbolic inputs
   e.g. angle = ang, speed = sp
- Collects symbolic constraints along each execution path e.g. consider 2 paths computing pressure2 in our example

## Symbolic Execution on our example

int calc\_trans\_slow\_torques (int angle, int speed)

```
{
  int pressure1, pressure2;
  int gear, val1, val3, val4;
  val1 = lookup1 (&(data_table[0][0]), angle);
  if (3 * speed \le val1)
    qear = 3;
  else
     gear = 4;
  val3 = lookup2 (\&(out1_table[0][0]), gear);
  pressure1 = val3 * 1000;
  val4 = lookup2 (\&(out2 table[0][0]), gear);
  pressure2 = val4 * 1000;
```

```
angle = 30, speed = 13
```

data table:	angle	val1
_	30	41

gear	val4
3	0
4	1

Path1 symbolic constraints:

out2\_table:

ang <=  $30 \land val1 = 41 \land 3^*$ sp <= val1  $\land$  gear =  $3 \land val4$ =  $0 \land pressure2 = val4^*$ 1000

9

## Symbolic Execution on our example

int calc\_trans\_slow\_torques (int angle, int speed)

```
int pressure1, pressure2;
```

int gear, val1, val3, val4;

```
val1 = lookup1 (&(data_table[0][0]), angle);
```

```
if (3 * speed <= val1)
```

gear = 3;

#### else

{

gear = 4;

```
val3 = lookup2 (&(out1_table[0][0]), gear);
pressure1 = val3 * 1000;
val4 = lookup2 (&(out2_table[0][0]), gear);
pressure2 = val4 * 1000;
```

angle = 
$$30$$
, speed =  $14$ 

data table:	angle	
	30	
out2_table:	gear	

gear	val4
3	0
4	1

val1

41

#### Path2 symbolic constraints:

ang <=  $30 \land val1 = 41 \land 3^*$ sp > val1  $\land$  gear =  $4 \land val4$ =  $1 \land pressure2 = val4^*$ 1000

## **Concolic Execution**

Symbolic testing technique

Concolic = Concrete + Symbolic



 Symbolic constraints for each explored path at the end of concolic execution

# Concolic Execution on our example

int calc\_trans\_slow\_torques (int angle, int speed)

{

int pressure1, pressure2;

int gear, val1, val3, val4;

```
val1 = lookup1 (&(data_table[0][0]), angle);
```

```
if (3 * speed \le val1)
```

gear = 3;

#### else

gear = 4;

val3 = lookup2 (&(out1\_table[0][0]), gear);

pressure1 = val3 \* 1000;

```
val4 = lookup2 (&(out2_table[0][0]), gear);
```

```
pressure2 = val4 * 1000;
```

Concrete input: angle = 30, speed = 13 Symbolic input: angle = ang, speed = sp

Symbolic path constraint: 3 \* sp <= val1

Negate the path conditional constraint:

Symbolic path constraint: 3 \* sp > val1

Solve the modified constraints:

New concrete input: angle = 30, speed = 14

Concrete input: angle = 30, speed = 14 Symbolic input: angle = ang, speed = sp

Symbolic path constraint: 3 \* sp > val1

## Symbolic Robustness Analysis

1. Path1 symbolic constraints:

ang <=  $30 \land val1 = 41 \land 3 * sp <= val1 \land$ gear =  $3 \land val4 = 0 \land pressure2 = val4 *$ 1000 2. Path2 symbolic constraints:

ang' <= 30 \langle val1' = 41 \langle 3 \* sp' > val1' \langle gear' = 4 \langle val4' = 1 \langle pressure2' = val4' \* 1000

 Formulate the optimization problem for above two paths: Maximize [pressure2 – pressure2']

subject to the constraints:

 Path1 symbolic constraints
 Path2 symbolic constraints angle = angle' |speed - speed'| <= 1</li>

 Iterate over all path pairs to find the maximum deviation in output (pressure2) for a perturbation of 1 unit (δ) in the input (speed).

Formal Problem Definition  

$$\epsilon_{yx} = \max_{v,x,x'} \begin{cases} |y - y'| & y = P(v,x) \\ y' = P(v,x') \\ |x - x'| \le \delta_x \end{cases}$$

- x : actual input value
- x' : measured input value
- v : value of all other input variables
- y : value of output of program P for input v, x
- y': value of output of program P for input v, x'
- $\delta_x$ : at most deviation of input variable x

 $\epsilon_{yx}$  :maximum deviation in the value of output variable y

Program P is (δ,ε)-robust in input x if  $ε_{vx} <= ε$  when  $δ_x <= δ$ 

## **Implementation Algorithm**

Algorithm find\_output\_sensitivity (P,  $x, \delta_x$ )



## Implementation

- Concolic execution: Splat
  - Uses the decision procedure STP to solve the path constraints to generate inputs for a new execution
  - Is extended to generate the symbolic output expression and the path constraints for all program paths
- Optimization problem solving: Lindo
  - Provides general nonlinear and nonlinear/integer optimization problem solving APIs
  - Can detect if a set of constraints is not satisfiable

## Limitations

- Limitations in algorithm and implementation:
  - Decision procedure handles fixed point numbers, not floating point numbers
  - $\delta_x$  is treated as constant
  - Closed loop systems can not be analyzed



In *n* iterations: input  $x_1$  will be  $2^n \delta$  apart

## Take Away

- Robustness analysis of control software program
- Symbolic technique of test case generation



- Robustness analysis of various types of entities:
  - Control Software programs
  - General Software programs
  - Sequential circuits
  - Networked systems
  - Cache replacement policies

## Idea

- Application of a approach used in robustness analysis of sequential circuits to analyze closed loop control software:
  - Sensor input variable values : actual and measured
  - Corresponding value sequences of state input variable
  - Corresponding value sequences of output variable
  - Common Suffix Distance:

last position of difference = loop iteration number

after which small change sensor input is forgotten

